

Parsing and Network Security

- Postel's law
- Parsing

Remotely Exploitable Vulnerabilities

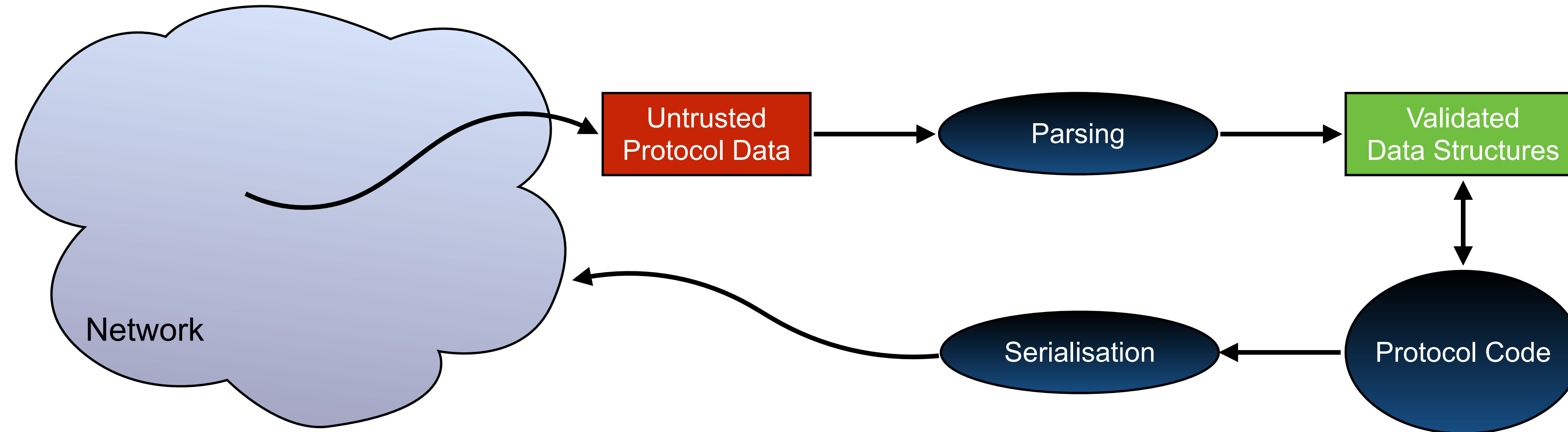
- Writing safe code in unsafe languages is difficult
- Easy to rationalise each individual bug – “How could anyone write that?”

```
if ((err = SSLHashSHA1.update(&hashCtx, &signedParams)) != 0)
    goto fail;
goto fail;
... other checks ...
fail:
    ... buffer frees (cleanups) ...
return err;
```

<https://dwheeler.com/essays/apple-goto-fail.html>

- But, people will continue to make mistakes
- Remotely exploitable vulnerabilities threaten any networked system
- **Are there particular classes of bug that cause such vulnerabilities?**

Parsing Untrusted Input



- The input parser is critical to security of a networked system
 - Takes untrusted input from the network
 - Generates validated, strongly typed, data structures that are processed by the rest of the code
- How can we ensure parsers are safe?

The Robustness Principle (Postel's Law)

At every layer of the protocols, there is a general rule whose application can lead to enormous benefits in robustness and interoperability:

`"Be liberal in what you accept, and conservative in what you send"`

Software should be written to deal with every conceivable error, no matter how unlikely; sooner or later a packet will come in with that particular combination of errors and attributes, and unless the software is prepared, chaos can ensue. In general, it is best to assume that the network is filled with malevolent entities that will send in packets designed to have the worst possible effect. This assumption will lead to suitable protective design, although the most serious problems in the Internet have been caused by un-envisaged mechanisms triggered by low-probability events; mere human malice would never have taken so devious a course!

Traditional guidance when writing networked systems is expressed in Postel's law

RFC1122

The Robustness Principle (Postel's Law)

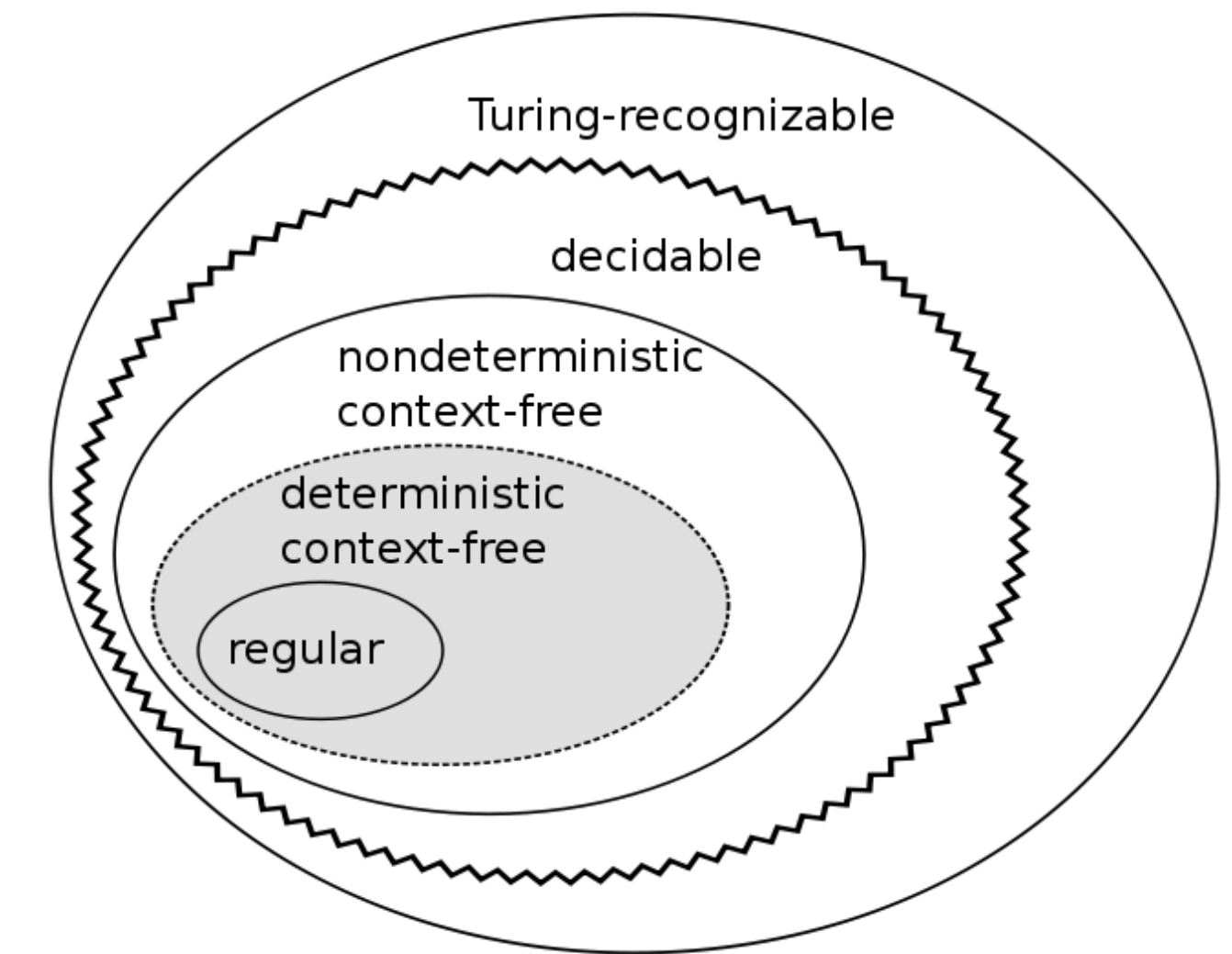
"Be liberal in what you accept, and conservative in what you send"

"Postel lived on a network with all his friends.
We live on a network with all our enemies.
Postel was wrong for today's internet."
— *Poul-Henning Kamp*

See also: <https://datatracker.ietf.org/doc/draft-iab-protocol-maintenance/>

Define Legal Inputs and Failure Responses

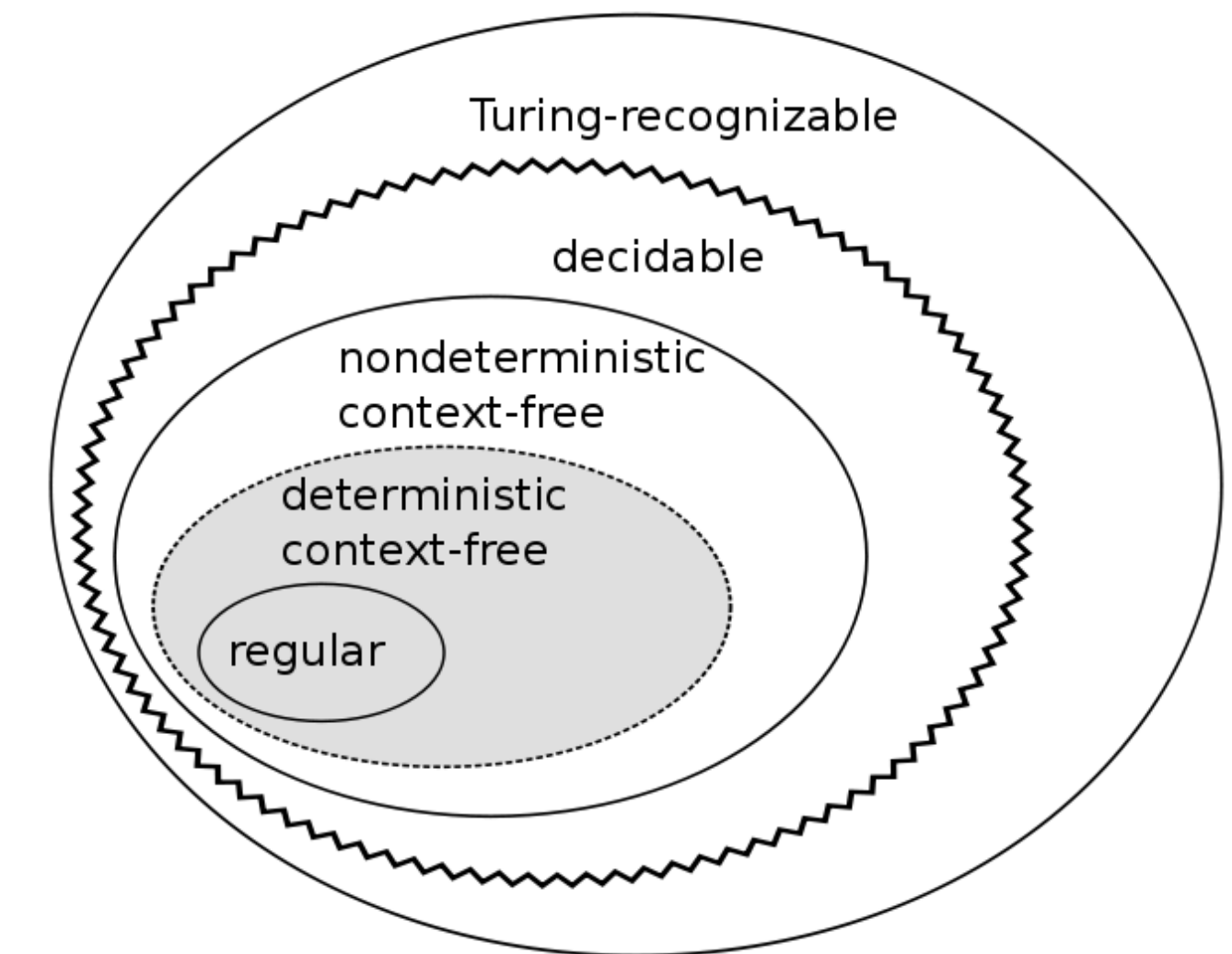
- The input parser takes untrusted input from the network, generates validated, strongly typed, data structures that are processed by the rest of the code
- A parser is an rule-driven automaton
 - It parses the input according to some grammar
 - If the input does not match the grammar, it fails
- **Formally specify the protocol grammar**
 - Define what is legal, and what is not, and write this down in a machine checkable manner
 - Define the grammar in as restrictive a manner as possible – e.g., don't use a Turing-complete parse when a regular expression will suffice
 - Specify what happens if the input data does not match the grammar: what causes a failure? How is it handled?



Source: <http://langsec.org/papers/Sassaman.pdf>

Generate the Parser (1/3)

- Difficult to reason about ad-hoc, manually written, parsing code
 - It performs low-level bit manipulation, string parsing, etc., all of which are hard to get correct
 - It tends to be poorly structured, hard to follow
- Rather, **auto-generate the parsing code:**
 - If input language is regular, use regular expression
 - If input language is context free, use a context free grammar
 - If neither of these work, use a more sophisticated parser, with minimal computational power
 - Generate strongly typed data structures, with explicit types to identify different data items
- Focus on parser correctness and readability
 - **Parsing performance matters less than security**



Generate the Parser (2/3)

- Use existing, well-tested, parser libraries
- For Rust code, use **nom** or **combine**:

nom, eating data byte by byte

<https://github.com/Geal/nom>

Combine: A parser combinator library for Rust

<https://github.com/Marwes/combine>

- For C or C++, use **hammer**

Hammer: Parser combinators for binary formats, in C. Yes, in C. What? Don't look at me like that.

<https://github.com/UpstandingHackers/hammer>



P. Chifflier and G. Couprie. Writing parsers like it is 2017. IEEE Workshop on Language-Theoretic Security, San Jose, CA, USA, May 2017. <https://dx.doi.org/10.1109/SPW.2017.39>

Generate the Parser (3/3)

```
# [derive (Debug, PartialEq, Eq) ]
pub struct Header {
    pub version: u8,
    pub audio:   bool,
    pub video:   bool,
    pub offset:  u32,
}

do_parse! (
    tag! ("FLV") >>
    version: be_u8 >>
    flags:   be_u8 >>
    offset:  be_u32 >>
    (Header {
        version: version,
        audio:   flags & 4 == 4,
        video:   flags & 1 == 1,
        offset:  offset
    })
)
);
```

- Specify the types into which parsed data is stored
- Describe the parser using an appropriate formal language – example uses **nom**
- Generate the parser from that language

- Parsing is performed first, and either succeeds in its entirety or fails – the rest of the code uses only safe, pre-parsed, data

Source: P. Chifflier and G. Couprie. Writing parsers like it is 2017. IEEE Workshop on Language-Theoretic Security, San Jose, CA, USA, May 2017. <https://dx.doi.org/10.1109/SPW.2017.39>

Define Parsable Protocols

- If designing network protocols, consider ease of parsing
 - Minimise the amount of state and look-ahead required to parse the data
 - Prefer a predictable, regular, grammar to one that saves bits at the expense of complex parsing
 - Networks get faster, security vulnerabilities remain
 - The benefit of saving a few bits gets less over time
 - The benefit of being easy to parse securely remains

Parsing and LangSec

- **Read:** “The Bugs We Have to Kill”
 - Is this approach realistic?
 - Can we combine better parsing with modern, strongly typed, languages to improve network security?
 - Is performance good enough?
 - Can we improve the way we design protocols?



The Bugs We Have to Kill
SERGEY BRATUS, MEREDITH L. PATTERSON, AND ANNA SHUBINA

Sergey Bratus is a Research Associate Professor of computer science at Dartmouth College. He sees state-of-the-art hacking as a distinct research and engineering discipline that, although not yet recognized as such, harbors deep insights into the nature of computing. He has a PhD in mathematics from Northeastern University and worked at BBN Technologies on natural language processing research before coming to Dartmouth.
sergey@cs.dartmouth.edu

Meredith L. Patterson is the founder of Upstanding Hackers. She developed the first language-theoretic defense against SQL injection in 2005 as a PhD student at the University of Iowa and has continued expanding the technique ever since. She lives in Brussels, Belgium.
mip@upstandinghackers.com

Anna Shubina is a Research Associate at the Dartmouth Institute for Security, Technology, and Society and maintains the CRAWDAD.org repository of traces and data for all kinds of wireless and sensor network research. She was the operator of Dartmouth's Tor node when the Tor network had about 30 nodes total.
ashubina@cs.dartmouth.edu

The code that parses inputs is the first and often the only protection for the rest of a program from malicious inputs. No programmer can afford to verify every implied condition on every line of code—even if this were possible to implement without slowing execution to a crawl. The parser is the part that is supposed to create a world for the rest of the program where all these implied conditions are true and need not be explicitly checked at every turn. Sadly, this is exactly where most parsers fail, and the rest of the program fails with them. In this article, we explain why parsers continue to be such a problem, as well as point to potential solutions that can kill large classes of bugs.

To do so, we are going to look at the problem from the computer science theory angle. Parsers, being input-consuming machines, are quite close to the theory's classic computing models, each one an input-consuming machine: finite automata, pushdown automata, and Turing machines. The latter is our principal model of general-purpose programming, the computing model with the ultimate power and flexibility. Yet this high-end power and flexibility come with a high price, which Alan Turing demonstrated (and to whose proof we owe our very model of general-purpose programming): our inability to predict, by any general static analysis algorithm, how programs for it will execute.

Yet most of our parsers are just a layer on top of this totally flexible computing model. It is not surprising, then, that without carefully limiting our parsers' design and code to much simpler models, we are left unable to prove these input-consuming machines secure. This is a powerful argument for making parsers and their input formats and protocols simpler, so that securing them does not require having to solve undecidable problems!

Parsers, Parsers Everywhere
To quote Karpowicz and Binszok [1]:

Parsing is of major interest in computer science. Classically discovered by students as the first step in compilation, parsing is present in almost every program which performs data-manipulation. For instance, the Web is built on parsers. The HyperText Transfer Protocol (HTTP) is a parsed dialog between the client, or browser, and the server. This protocol transfers pages in HyperText Markup Language (HTML), which is also parsed by the browser. When running web-applications, browsers interpret JavaScript programs which, again, begins with parsing. Data exchange between browser(s) and server(s) uses languages or formats like XML and JSON. Even inside the server, several components (for instance the trio made of the HTTP server Apache, the PHP interpreter and the MySQL database) often manipulate programs and data dynamically; all require parsers.

So do the lower layers of the network stack down to the IP and the link layer protocols, and also other OS parts such as the USB drivers [2] (and even the hardware: turning PHY layer symbol streams into frames is parsing, too!). For all of these core protocols, we add, their parsers have had a long history of failures, resulting in an Internet where any site, program, or system that receives untrusted input can be presumed compromised.

4 **login:** AUGUST 2015 VOL. 40, NO. 4 www.usenix.org

S. Bratus, M. L. Patterson, and A. Shubina. The bugs we have to kill. *login:*, 40(4):4–10, August 2015. <http://langsec.org/papers/the-bugs-we-have-to-kill.pdf>

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