



University  
of Glasgow

# Network Layer (2)

Networked Systems 3  
Lecture 9

# Lecture Outline

- Addressing
  - Concepts
  - Addressing in the Internet
    - IPv4
    - IPv6
    - The shortage of IPv4 addresses

# Addressing

- How to name hosts in a network?
  - Is the address an identity or a location?
    - Does it name the host, or the location at which it attaches to the network
  - How should addresses be allocated?
    - Hierarchical or flat?
  - What is the address format?
    - Human or machine readable?
    - Textual or binary? Structured or unstructured?
    - Fixed or variable length? How large?


# Identity and Location

- Addresses can denote host identity
  - Give hosts a consistent address, irrespective of where or when they attach to the network
  - Simple upper-layer protocols
    - Transport layer and applications unaware of multi-homing or host mobility
  - Puts complexity in network layer
    - Network must determine location of host before it can route data
    - Often requires in-network database to map host identity to routable address
    - E.g. non-geographic and mobile phone numbers: 0870 154 154 → 01604 230 230

# Identity and Location

- Alternatively, an address can indicate the *location* at which a host attaches to the network
  - Address structure matches the network structure
    - Network can directly route data given an address
    - E.g. geographic phone numbers: +44 141 330 4256
  - Simplifies network layer, by pushing complexity to the higher layers
    - Multi-homing and mobility must be handled by transport layer or applications
    - E.g. transport layer connections break when host moves

# Address Allocation

- Are addresses allocated hierarchically?
  - Allows routing on aggregate addresses
    - E.g. phone call to +1 703 243 9422  
 Route to US without looking at rest of number
  - Forces address structure to match network topology
  - Requires rigid control of allocations
- Or is there a flat namespace?
  - Flexible allocations, no aggregation → not scalable

# Address Formats

- Textual or binary? Fixed or variable length?
  - Fixed length binary easier (faster) for machines to process
  - Variable length textual easier for humans to read
  - Which are you optimising for?

# IP Addresses

- IP addresses have the following characteristics:
  - They specify location of a network interface
  - They are allocated hierarchically
  - They are fixed length binary values
    - IPv4: 32 bits
    - IPv6: 128 bits
- Domain names are a separate *application level* namespace



# IP Addresses

- Both IPv4 and IPv6 addresses encode location
  - Addresses are split into a *network part* and a *host part*
    - A *netmask* describes the number of bits in the network part
    - The network itself has the address with the host part equal to zero
    - The broadcast address for a network has all bits of host part equal to one(allows messages to be sent to all hosts on a network)
  - A host with several network interfaces will have one IP addresses per interface
    - E.g. a laptop with Ethernet and Wi-Fi interfaces will have two IP addresses

# IPv4 Addresses

## 32 bit binary addresses

IP address: 130.209.247.112 = 10000010 11010001 11110111 01110000

Netmask: 255.255.240.0 = 11111111 11111111 11110000 00000000



20 bits → network = 130.209.240.0/20

Broadcast address:

130.209.255.255 = 10000010 11010001 11111111 11111111

# Aside: Classes of IP address

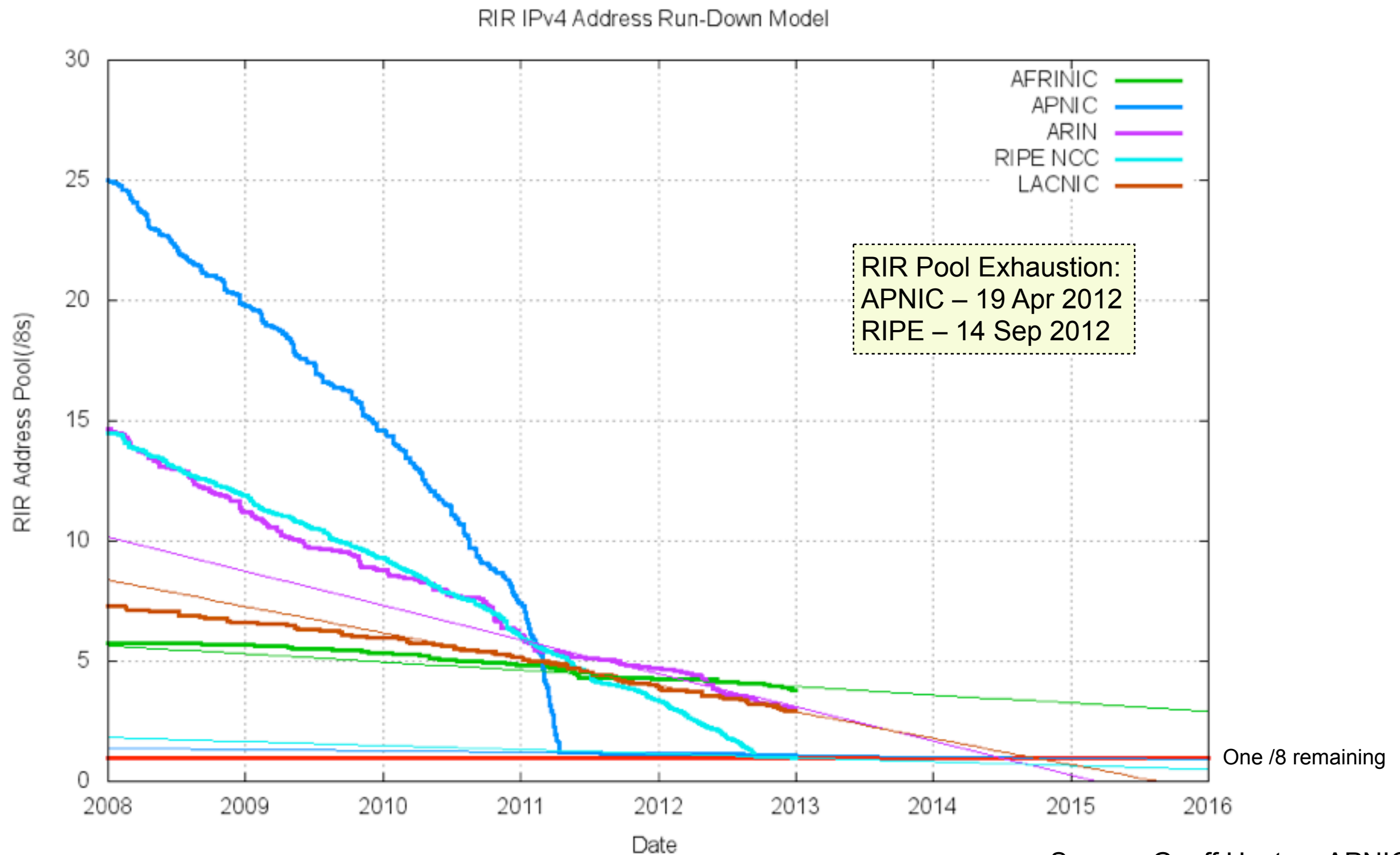
- IP addresses used to be allocated so the netmask was a multiple of 8 bits
  - Class A → a /8 network (~16 million addresses)
  - Class B → a /16 network (65536 addresses)
  - Class C → a /24 network (256 addresses)
- Inflexible, and wasted addresses
- Arbitrary length netmask allowed since 1993:
  - The DCS network is a /20

Old terminology still used sometimes...

# IP Address Management

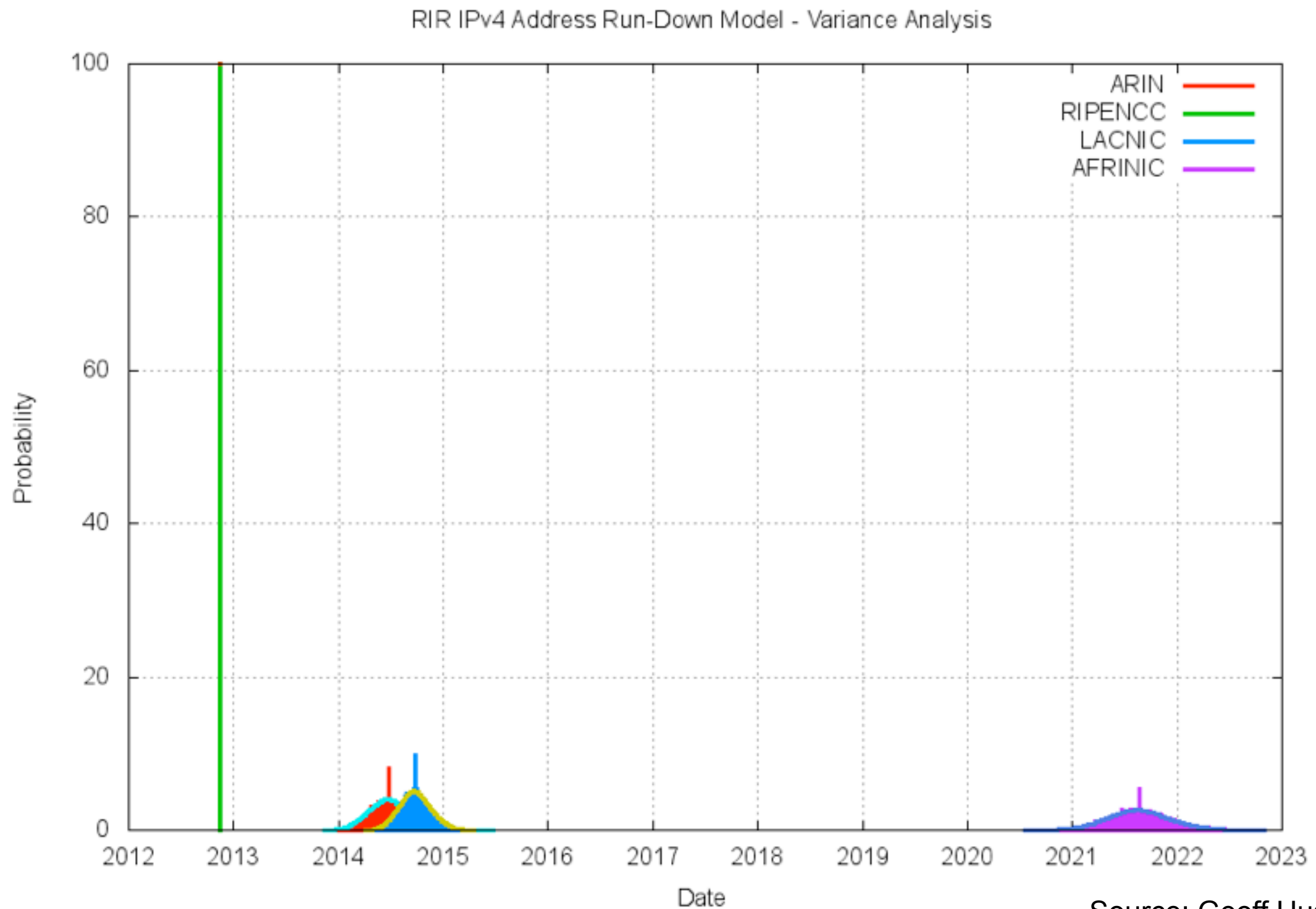
- IPv4 has  $2^{32} = 4,294,967,296$  addresses
  - IANA administers the pool of unallocated addresses
    - Historically would assign addresses directly to ISPs, large enterprises, etc.
    - Now, addresses assigned to regional Internet registries (RIRs) as needed:
      - AfriNIC (Africa), APNIC (Asia-Pacific), ARIN (North America), LACNIC (Latin America and Caribbean), and RIPE (Europe, Middle East, Central Asia)
      - Allocations made one /8 ( $2^{24} = 16,777,216$  addresses) at a time
  - RIRs allocate addresses to ISPs and large enterprises within their region; ISPs allocate to their customers
- IANA has allocated all available addresses to RIRs
  - Last allocation made on 3 February 2011

# RIR Exhaustion Dates (1)



Source: Geoff Huston, APNIC  
2 Jan 2013 <http://ipv4.potaroo.net/>

# RIR Exhaustion Dates (2)



Source: Geoff Huston, APNIC  
2 Jan 2013 <http://ipv4.potaroo.net/>

# What Happens Next?

- IPv6 is widely deployed
- Widespread use of NAT
- Or...?





# IPv6

- IPv6 provides 128 bit addresses – if deployed it will solve address shortage for a *long* time
- $2^{128} =$   
340,282,366,920,938,463,463,374,607,431,768,211,456  
addresses
- i.e., 665,570,793,348,866,943,898,599 addresses  
for every square metre of the Earth's surface
- i.e., the same number of IP addresses per person  
as the number of atoms in a metric ton of carbon



# IPv6 Addresses

128 bit binary addresses, written as : separated hexadecimal

2001:0db8:85a3:08d3:1319:8a2e:0370:7334

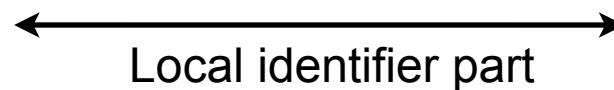
*A single* run of consecutive zeros can be compressed to a ::

2001:200::8002:203:47ff:fea5:3085

# IPv6 Addresses

Local identifier part of IPv6 address is 64 bits:

2001:0db8:85a3:08d3:1319:8a2e:0370:7334



Can be derived from Ethernet/Wi-Fi MAC address:

48 bit IEEE MAC: 0014:5104:25ea

Expand to 64 bits: 0014:51ff:fe04:25ea

Invert bit 6: 0214:51ff:fe04:25ea

Or randomly chosen, with bit 6 set  
to zero, to give illusion of privacy

# IPv6 Addresses

Routers advertise network part, hosts auto-configure address:

2001:0db8:85a3:08d3:1319:8a2e:0370:7334

← Network part →

Network part is split into a *global routing prefix* (a.k.a. “routing goop” of up to 48 bits) and a *subnet identifier*

Formalises the distinction present in IPv4:

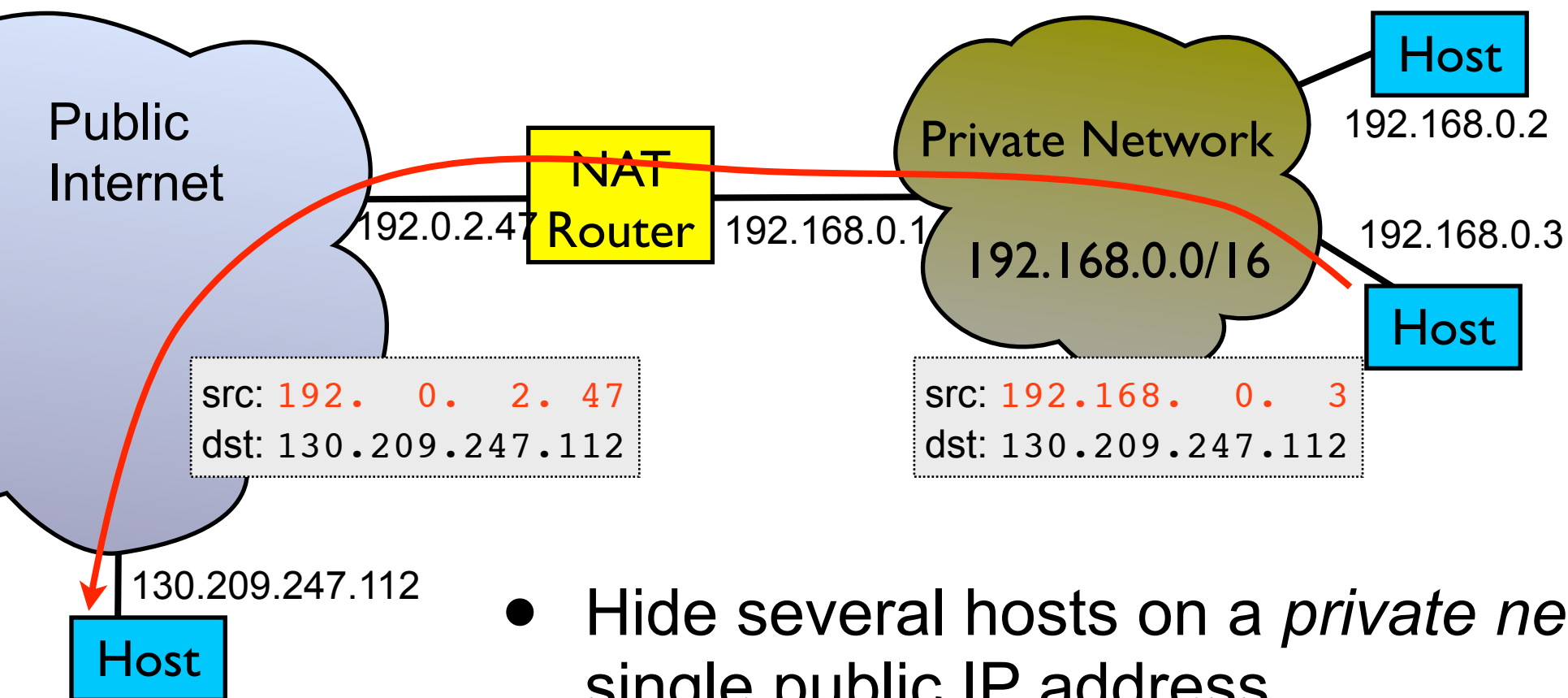
130.209.247.112 = 10000010 11010001 11110111 01110000

← Global routing prefix Subnet Host →  
← Netmask →

# IPv6 Deployment Issues

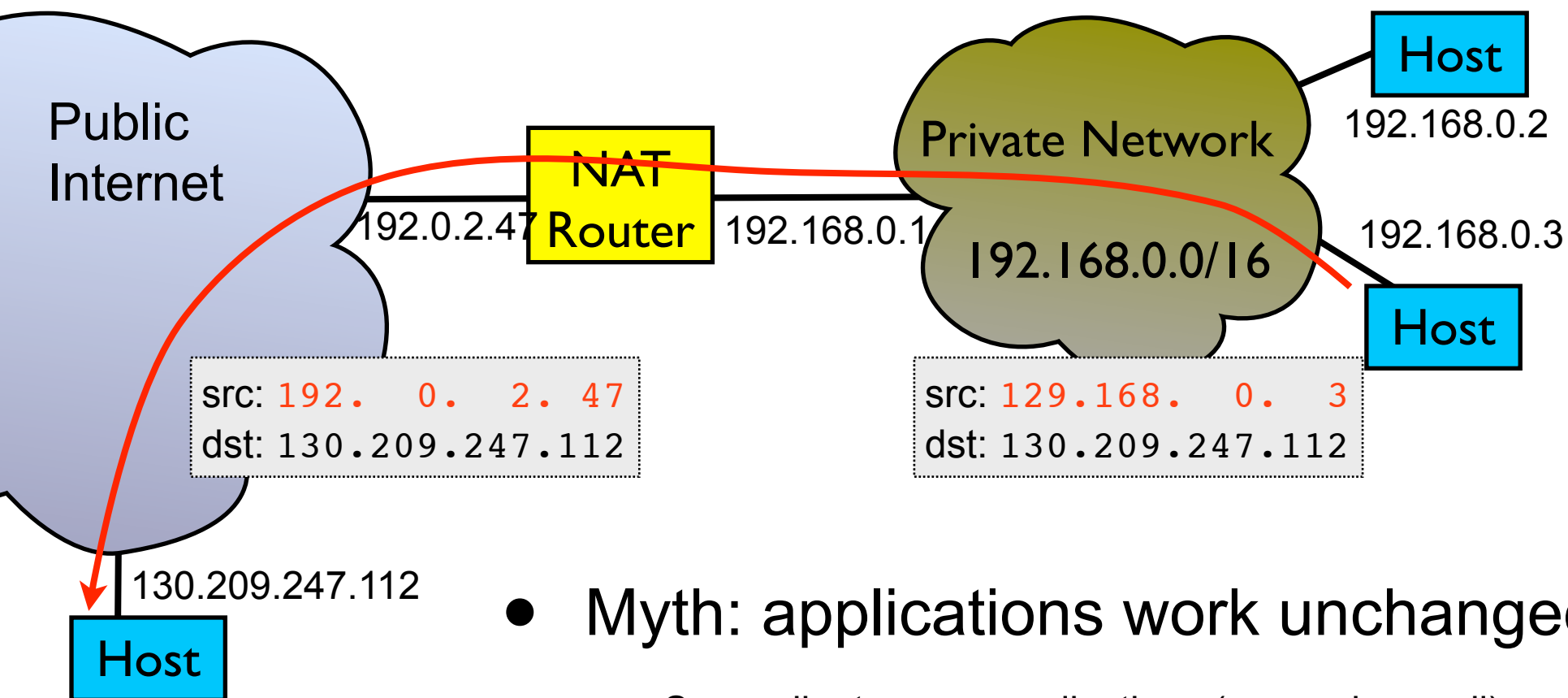
- IPv6 requires changes to *every* single host, router, firewall, and application...
  - Significant deployment challenge!
  - Host changes done: MacOS X, Windows XP & Vista, Linux, FreeBSD, Symbian, etc.
  - Backbone routers generally support IPv6, home routers and firewalls (mostly) not yet
  - Some applications have been updated

# Network Address Translation



- Hide several hosts on a *private network* behind a single public IP address
  - Private IPv4 addresses are 10.0.0.0/8, 192.168.0.0/16, 176.16.0.0/12
- Rewrite packet headers at network boundary
  - Doesn't require changes to hosts or routers (other than the NAT)
- Tries to give the illusion of more address space

# Network Address Translation



- **Myth: applications work unchanged**
  - Some client-server applications (e.g. web, email) work without changes
  - But peer-to-peer applications (e.g. VoIP) need extensive changes before they work through a NAT (~200 pages spec to describe algorithm!)
- **Myth: provides security**
  - Most NATs include a firewall to provide security, the NAT function gives no security benefit

# NAT vs. IPv6

- NAT widely deployed now for IPv4
  - Initially *appears* simple to end users (no need to change hosts, web browsing & email still work)
  - *Hugely* complex for peer-to-peer applications
  - Very difficult to debug network problems and deploy new classes of application
- The *hope* is that IPv6 will be deployed as IPv4 addresses run out

# Questions?