RTP Control Protocol (RTCP) Feedback for Congestion Control
draft-dt-rmcat-feedback-message-03

Abstract

This document describes a feedback message intended to enable congestion control for interactive real-time traffic. The RTP Media Congestion Avoidance Techniques (RMCAT) Working Group formed a design team to analyze feedback requirements from various congestion control algorithms and to design a generic feedback message to help ensure interoperability across those algorithms. The feedback message is designed for a sender-based congestion control, which means the receiver of the media will send necessary feedback to the sender of the media to perform the congestion control at the sender.

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1. Introduction

For interactive real-time traffic the typical protocol choice is Realtime Transport Protocol (RTP) over User Datagram Protocol (UDP). RTP does not provide any guarantee of Quality of Service (QoS), reliable or timely delivery and expects the underlying transport protocol to do so. UDP alone certainly does not meet that expectation. However, RTP Control Protocol (RTCP) provides a mechanism to periodically send transport and media metrics to the media sender which can be utilized and extended for the purposes of RMCAT congestion control. For a congestion control algorithm which operates at the media sender, RTCP messages can be transmitted from the media receiver back to the media sender to enable congestion control. In the absence of standardized messages for this purpose, the congestion control algorithm designers have designed proprietary RTCP messages that convey only those parameters required for their respective designs. As a direct result, the different congestion control (a.k.a. rate adaptation) designs are not interoperable. To enable algorithm evolution as well as interoperability across designs...
(e.g., different rate adaptation algorithms), it is highly desirable to have generic congestion control feedback format.

To help achieve interoperability for unicast RTP congestion control, this memo proposes a common RTCP feedback format that can be used by NADA [I-D.ietf-rmcat-nada], SCReAM [I-D.ietf-rmcat-scream-cc], Google Congestion Control [I-D.ietf-rmcat-gcc] and Shared Bottleneck Detection [I-D.ietf-rmcat-sbd], and hopefully future RTP congestion control algorithms as well.

2. Terminology

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in [RFC2119].

In addition the terminology defined in [RFC3550], [RFC3551], [RFC3611], [RFC4585], and [RFC5506] applies.

3. Feedback Message

The design team analyzed the feedback requirements from the different proposed candidate in RMCAT WG. The analysis showed some commonalities between the proposed solution candidate and some can be derived from other information. The design team has agreed to have following packet information block in the feedback message to satisfy different requirement analyzed.

- **Packet Identifier**: RTP sequence number. The RTP packet header includes an incremental packet sequence number that the sender needs to correlate packets sent at the sender with packets received at the receiver.

- **Packet Arrival Time**: Arrival time stamp at the receiver of the media. The sender requires the arrival time stamp of the respective packet to determine delay and jitter the packet had experienced during transmission. In a sender based congestion control solution the sender requires to keep track of the sent packets – usually packet sequence number, packet size and packet send time. With the packet arrival time the sender can detect the delay and jitter information. Along with packet loss and delay information the sender can estimate the available bandwidth and thus adapt to the situation.

- **Packet Explicit Congestion Notification (ECN) Marking**: If ECN [RFC3168] is used, it is necessary to report on the 2-bit ECN mark in received packets, indicating for each packet whether it is marked not-ECT, ECT(0), ECT(1), or ECN-CE. If the path on which
the media traffic traversing is ECN capable then the sender can use the Congestion Experienced (ECN-CE) marking information for congestion control. It is important that the receiver sends the ECN-CE marking information of the packet back to the sender to take the advantages of ECN marking. Note that how the receiver gets the ECN marking information at application layer is out of the scope of this design team. Additional information for ECN use with RTP can be found at [RFC6679].

The feedback messages can have one or more of the above information blocks. For RTCP based feedback message the packet information block will be grouped by Synchronization Source (SSRC) identifier.

As a practical matter, we note that host Operating System (OS) process interruptions can occur at inopportune times. Thus, the recording of the sent times at the sender and arrival times at the receiver should be made with deliberate care. This is because the time duration of host OS interruptions can be significant relative to the precision desired in the one-way delay estimates. Specifically, the send time should be recorded at the latest opportunity prior to outputting the media packet at the sender (e.g., socket or RTP API) and the arrival time at the receiver (e.g., socket or RTP API) should be recorded at the earliest opportunity available to the receiver.

3.1. RTCP XR Block for Reporting Congestion Control Feedback

Congestion control feedback can be sent as part of a scheduled RTCP report, or as RTP/AVPF early feedback. If sent as part of a scheduled RTCP report, the feedback is sent as an XR block, as part of a regular compound RTCP packet. The format of the RTCP XR report block is as follows (this will be preceded in the compound RTCP packet by an RTCP XR header, described in [RFC3611], that includes the SSRC of the report sender):
The XR Discard RLE report block uses the same format as specified for the loss and duplicate report blocks in [RFC3611]. The fields "block length", "begin_seq", and "end_seq" have the same semantics and representation as defined in [RFC3611].

Block Type (BT, 8 bits): The RMCAT congestion control feedback Report Block is identified by the constant RC2F. [Note to RFC Editor: Please replace RC2F with the IANA provided RTCP XR block type for this block.]

Report Count (8 bits): field describes the number of SSRCs reported by this report block. The number should at least be 1.

Report Timestamp (RTS, 32 bits): represents the timestamp when this report was generated. The sender of the feedback message decides on the wall-clock. Usually, it should be derived from the same wall-clock that is used for timestamping RTP packets arrival. Consistency in the unit and resolution (10th of millisecond should be good enough) is important here. In addition, the media sender can ask for a specific resolution it wants.
Each sequence number between the begin_seq and end_seq (both inclusive) is represented by a packet metric block of 16-bits that contains the L, ECN, and ATO metrics. If an odd number of reports are included, i.e., end_seq - begin_seq is odd then 16 bits of zero padding MUST be added after the last report, to align the RTCP packet to a four (4) bytes boundary.

L (1 bit): is a boolean to indicate if the packet was received. 0 represents that the packet was not yet received and all the subsequent bits (ECN and ATO) are also set to 0. 1 represent the packet was received and the subsequent bits in the block need to be parsed.

ECN (2 bits): is the echoed ECN mark of the packet (00 if not received or if ECN is not used).

Arrival time offset (ATO, 13 bits): it the relative arrival time of the RTP packets at the receiver before this feedback report was generated measured in milliseconds. It is calculated by subtracting the reception timestamp of the RTP packet denoted by this 16bit block and the timestamp (RTS) of this report. If the measured value is greater than 8.189 seconds (the value that would be coded as 0x1FFD), the value 0x1FFE MUST be reported to indicate an over-range positive measurement. If the measurement is unavailable, the value 0x1FFF MUST be reported.

3.2. RTP/AVPF Transport Layer Feedback for Congestion Control

Congestion control feedback can also be sent in a non-compound RTCP packet [RFC5506] if the RTP/AVPF profile [RFC4585] or the RTP/SAVPF profile [RFC5124] is used. In this case, the congestion control feedback is sent as a Transport Layer FB message (RTCP packet type 205), with FMT=2 (congestion control feedback). The format of this RTCP packet is as follows:
The first 8 octets are the RTCP header, with PT=205 and FMT=2 specifying the remainder is a congestion control feedback packet, and including the SSRC of the packet sender.

Section 6.1 of [RFC4585] requires this is followed by the SSRC of the media source being reported upon. Accordingly, the format of the report is changed from that of the RTCP XR report block, to move the timestamp to the end. The meaning of all the fields is described in Section 3.1.

4. Feedback Frequency and Overhead

There is a trade-off between speed and accuracy of reporting, and the overhead of the reports. [I-D.ietf-rmcat-rtp-cc-feedback] discusses this trade-off, and the possible rates of feedback.

It is a general understanding that the congestion control algorithms will work better with more frequent feedback – per packet feedback. However, RTCP bandwidth and transmission rules put some upper limits on how frequently the RTCP feedback messages can be send from the
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media receiver to the media sender. It has been shown [I-D.ietf-rmcat-rtp-cc-feedback] that in most cases a per frame feedback is a reasonable assumption on how frequent the RTCP feedback messages can be transmitted. The design team also have noted that even if a higher frequency of feedback is desired it is not viable if the feedback messages starts to compete against the media traffic on the feedback path during congestion period. Analyzing the feedback interval requirement [feedback-requirements] it can be seen that the candidate algorithms can perform with a feedback interval range of 50-200ms. A value within this range need to be negotiated at session setup.

5. Design Rationale

The primary function of RTCP Sender Report (SR) / Receiver Report (RR) is to report the reception quality of media. The regular SR / RR reports contain information about observed jitter, fractional packet loss and cumulative packet loss. The original intent of this information was to assist flow and congestion control mechanisms. Even though it is possible to do congestion control based on information provided in the SR/RR reports it is not sufficient to design an efficient congestion control algorithm for interactive real-time communication. An efficient congestion control algorithm requires more fine grain information on per packet (see Section 3) to react to the congestion or to avoid further congestion on the path.

Codec Control Message for AVPF [RFC5104] defines Temporary Maximum Media Bit Rate (TMMBR) message which conveys a temporary maximum bitrate limitation from the receiver of the media to the sender of the media. Even though it is not designed to replace congestion control, TMMBR has been used as a means to do receiver based congestion control where the session bandwidth is high enough to send frequent TMMBR messages especially with reduced sized reports [RFC5506]. This requires the receiver of the media to analyze the data reception, detect congestion level and recommend a maximum bitrate suitable for current available bandwidth on the path with an assumption that the sender of the media always honors the TMMBR message. This requirement is completely opposite of the sender based congestion control approach. Hence, TMMBR cannot be as a signaling means for a sender based congestion control mechanism. However, TMMBR should be viewed a complimentary signaling mechanism to establish receiver’s current view of acceptable maximum bitrate which a sender based congestion control should honor.

There are number of RTCP eXtended Report (XR) blocks have been defined for reporting of delay, loss and ECN marking. It is possible to combine several XR blocks to report the loss and ECN marking at
the cost of overhead and complexity. However, there is no existing RTCP XR block to report packet arrival time.

Considering the issues discussed here it is rational to design a new congestion control feedback signaling mechanism for sender based congestion control algorithm.

6. Acknowledgements

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7. IANA Considerations

7.1. RTP/AVPF Transport Layer Feedback Message

TBD

7.2. RTCP XR Report Blocks

TBD

8. Security Considerations

There is a risk of causing congestion if an on-path attacker modifies the feedback messages in such a manner to make available bandwidth greater than it is in reality. [More on security consideration TBD.]

9. References

9.1. Normative References

[I-D.ietf-rmcat-rtp-cc-feedback]

9.2. Informative References

[feedback-requirements]
"RMCAT Feedback Requirements",


Authors’ Addresses

Zaheduzzaman Sarker
Ericsson AB
Luleae
Sweden

Phone: +46107173743
Email: zaheduzzaman.sarker@ericsson.com

Colin Perkins
University of Glasgow
School of Computing Science
Glasgow G12 8QQ
United Kingdom

Email: csp@csperkins.org
Varun Singh
Nemu Dialogue Systems Oy
Runeberginkatu 4c A 4
Helsinki 00100
Finland

Email: varun.singh@iki.fi
URI: http://www.callstats.io/

Michael A. Ramalho
Cisco Systems, Inc.
6310 Watercrest Way Unit 203
Lakewood Ranch, FL 34202
USA

Phone: +1 919 476 2038
Email: mramalho@cisco.com