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RTP and the Datagram Congestion Control Protocol (DCCP)
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Abstract

The Real-time Transport Protocol (RTP) is a widely used transport for real-time multimedia on IP networks. The Datagram Congestion Control Protocol (DCCP) is a newly defined transport protocol that provides desirable services for real-time applications. This memo specifies a mapping of RTP onto DCCP, along with associated signalling, such that real-time applications can make use of the services provided by DCCP.

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1. Introduction

The Real-time Transport Protocol (RTP) [1] is widely used in video streaming, telephony, and other real-time networked applications. RTP can run over a range of lower-layer transport protocols, and the performance of an application using RTP is heavily influenced by the choice of lower-layer transport. The Datagram Congestion Control Protocol (DCCP) [2] is a newly specified transport protocol that provides desirable properties for real-time applications running on unmanaged best-effort IP networks. This memo describes how RTP can be framed for transport using DCCP, and discusses some of the implications of such a framing. It also describes how the Session Description Protocol (SDP) [3] can be used to signal such sessions.

The remainder of this memo is structured as follows: it begins with a rationale for the work in Section 2, describing why a mapping of RTP onto DCCP is needed. Following a description of the conventions used in this memo in Section 3, the specification begins in Section 4 with the definition of how RTP packets are framed within DCCP. Associated signalling is described in Section 5. Security considerations are discussed in Section 6, and IANA considerations in Section 7.

2. Rationale

With the widespread adoption of RTP have come concerns that many real time applications do not implement congestion control, leading to the potential for congestion collapse of the network [15]. The designers of RTP recognised this issue, stating that [4]:

If best-effort service is being used, RTP receivers SHOULD monitor packet loss to ensure that the packet loss rate is within acceptable parameters. Packet loss is considered acceptable if a TCP flow across the same network path and experiencing the same network conditions would achieve an average throughput, measured on a reasonable time-scale, that is not less than the RTP flow is achieving. This condition can be satisfied by implementing congestion control mechanisms to adapt the transmission rate (or the number of layers subscribed for a layered multicast session), or by arranging for a receiver to leave the session if the loss rate is unacceptably high.

While the goals are clear, the development of TCP friendly congestion control that can be used with RTP and real-time media applications is an open research question with many proposals for new algorithms, but little deployment experience.

Two approaches have been used to provide congestion control for RTP:

1) develop RTP extensions that incorporate congestion control; and 2) provide mechanisms for running RTP over congestion controlled transport protocols. An example of the first approach can be found in [16], extending RTP to incorporate feedback information such that TFRC congestion control [17] can be implemented at the application level. This will allow congestion control to be added to existing applications without operating system or network support, and it offers the flexibility to experiment with new congestion control algorithms as they are developed. Unfortunately, it also passes the complexity of implementing congestion control onto application authors, a burden which many would prefer to avoid.

The second approach is to run RTP on a lower-layer transport protocol that provides congestion control. One possibility is to run RTP over TCP, as defined in [5], but the reliable nature of TCP and the dynamics of its congestion control algorithm make this inappropriate for most interactive real time applications (the Stream Control Transmission Protocol (SCTP) is inappropriate for similar reasons). A better fit for such applications may be to run RTP over DCCP, since DCCP offers unreliable packet delivery and a choice of congestion control. This gives applications the ability to tailor the transport to their needs, taking advantage of better congestion control algorithms as they come available, while passing complexity of implementation to the operating system. If DCCP should come to be widely available, it is believed these will be compelling advantages. Accordingly, this memo defines a mapping of RTP onto DCCP.

3. Conventions Used in this Memo

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in RFC 2119 [6].

4. RTP over DCCP: Framing

The following section defines how RTP and RTCP packets can be framed for transport using DCCP. It also describes the differences between RTP sessions and DCCP connections, and the impact these have on the design of applications.

4.1. RTP Data Packets

Each RTP data packet MUST be conveyed in a single DCCP datagram. Fields in the RTP header MUST be interpreted according to the RTP specification, and any applicable RTP Profile and Payload Format. Header processing is not affected by DCCP framing (in particular,

note that the semantics of the RTP sequence number and the DCCP sequence number are not compatible, and the value of one cannot be inferred from the other).

A DCCP connection is opened when an end system joins an RTP session, and it remains open for the duration of the session. To ensure NAT bindings are kept open, an end system SHOULD send a zero length DCCP-Data packet once every 15 seconds during periods when it has no other data to send. This removes the need for RTP no-op packets [18], and similar application level keep-alives, when using RTP over DCCP. This application level keepalive does not need to be sent if it is known that the DCCP CCID in use provides a transport level keepalive.

RTP data packets MUST obey the dictates of DCCP congestion control. In some cases, the congestion control will require a sender to send at a rate below that which the payload format would otherwise use. To support this, an application could use either a rate adaptive payload format, or a range of payload formats (allowing it to switch to a lower rate format if necessary). Details of the rate adaptation policy for particular payload formats are outside the scope of this memo (but see [19] and [20] for guidance).

RTP extensions that provide application-level congestion control (e.g. [16]) will conflict with DCCP congestion control, and MUST NOT be used.

DCCP allows an application to choose the checksum coverage, using a partial checksum to allow an application to receive packets with corrupt payloads. Some RTP Payload Formats (e.g. [21]) can make use of this feature in conjunction with payload-specific mechanisms to improve performance when operating in environments with frequent non-congestive packet corruption. If such a payload format is used, an RTP end system MAY enable partial checksums at the DCCP layer, in which case the checksum MUST cover at least the DCCP and RTP headers to ensure packets are correctly delivered. Partial checksums MUST NOT be used unless supported by mechanisms in the RTP payload format.

4.2. RTP Control Packets

The RTP Control Protocol (RTCP) is used in the standard manner with DCCP. RTCP packets are grouped into compound packets, as described in Section 6.1 of [1], and each compound RTCP packet is transported in a single DCCP datagram.

The usual RTCP timing rules apply, with the additional constraint that RTCP packets MUST obey the DCCP congestion control algorithm negotiated for the connection. This can prevent a participant from sending an RTCP packet at the expiration of the RTCP transmission

timer if there is insufficient network capacity available. In such cases the RTCP packet is delayed and sent at the earliest possible instant when capacity becomes available. The actual time the RTCP packet was sent is then used as the basis for calculating the next RTCP transmission time.

RTCP packets comprise only a small fraction of the total traffic in an RTP session. Accordingly, it is expected that delays in their transmission due to congestion control will not be common, provided the configured nominal "session bandwidth" (see Section 6.2 of [1]) is in line with the bandwidth achievable on the DCCP connection. If, however, the capacity of the DCCP connection is significantly below the nominal session bandwidth, RTCP packets may be delayed enough for participants to time out due to apparent inactivity. In such cases, the session parameters SHOULD be re-negotiated to more closely match the available capacity, for example by performing a re-invite with an updated "b=" line when using the Session Initiation Protocol [22] for signalling.

Note: Since the nominal session bandwidth is chosen based on media codec capabilities, a session where the nominal bandwidth is much larger than the available bandwidth will likely become unusable due to constraints on the media channel, and so require negotiation of a lower bandwidth codec, before it becomes unusable due to constraints on the RTCP channel.

As noted in Section 17.1 of [2], there is the potential for overlap between information conveyed in RTCP packets and that conveyed in DCCP acknowledgement options. In general this is not an issue since RTCP packets contain media-specific data that is not present in DCCP acknowledgement options, and DCCP options contain network-level data that is not present in RTCP. Indeed, there is no overlap between the five RTCP packet types defined in the RTP specification [1] and the standard DCCP options [2]. There are, however, cases where overlap does occur: most clearly between the optional RTCP Extended Reports Loss RLE Blocks [23] and the DCCP Ack Vector option. If there is overlap between RTCP report packets and DCCP acknowledgements, an application SHOULD use either RTCP feedback or DCCP acknowledgements, but not both (use of both types of feedback will waste available network capacity, but is not otherwise harmful).

4.3. Multiplexing Data and Control

The obvious mapping of RTP onto DCCP creates two DCCP connections for each RTP flow: one for RTP data packets, one for RTP control packets. A frequent criticism of RTP relates to the number of ports it uses, since large telephony gateways can support more than 32768 RTP flows between pairs of gateways, and so run out of UDP ports. In addition,

use of multiple ports complicates NAT traversal. For these reasons, it is RECOMMENDED that the RTP and RTCP traffic for a single RTP session is multiplexed onto a single DCCP connection following the guidelines in [7], where possible (it may not be possible in all circumstances, for example when translating from an RTP stream over a non-DCCP transport that uses conflicting RTP payload types and RTCP packet types).

4.4. RTP Sessions and DCCP Connections

An end system SHOULD NOT assume that it will observe only a single RTP synchronisation source (SSRC) because it is using DCCP framing. An RTP session can span any number of transport connections, and can include RTP mixers or translators bringing other participants into the session. The use of a unicast DCCP connection does not imply that the RTP session will have only two participants, and RTP end systems SHOULD assume that multiple synchronisation sources may be observed when using RTP over DCCP, unless otherwise signalled.

An RTP translator bridging multiple DCCP connections to form a single RTP session needs to be aware of the congestion state of each DCCP connection, and must adapt the media to the available capacity of each. The Codec Control Messages defined in [24] may be used to signal congestion state to the media senders, allowing them to adapt their transmission. Alternatively, media transcoding may be used to perform adaptation: this is computationally expensive, induces delay, and generally gives poor quality results. Depending on the payload, it might be possible to use some form of scalable coding. Scalable media coding formats are an active research area, and are not in widespread use at the time of this writing.

A single RTP session may also span a DCCP connection and some other type of transport connection. An example might be an RTP over DCCP connection from an RTP end system to an RTP translator, with an RTP over UDP/IP multicast group on the other side of the translator. A second example might be an RTP over DCCP connection that links PSTN gateways. The issues for such an RTP translator are similar to those when linking two DCCP connections, except that the congestion control algorithms on either side of the translator may not be compatible. Implementation of effective translators for such an environment is non-trivial.

4.5. RTP Profiles

In general, there is no conflict between new RTP Profiles and DCCP framing, and most RTP profiles can be negotiated for use over DCCP with the following exceptions:

- o An RTP profile that is intolerant of packet corruption may conflict with the DCCP partial checksum feature. An example of this is the integrity protection provided by the RTP/SAVP profile, which cannot be used in conjunction with DCCP partial checksums.
- o An RTP profile that mandates a particular non-DCCP lower layer transport will conflict with DCCP.

RTP profiles which fall under these exceptions SHOULD NOT be used with DCCP unless the conflicting features can be disabled.

Of the profiles currently defined, the RTP Profile for Audio and Video Conferences with Minimal Control [4], the Secure Real-time Transport Protocol [8], the Extended RTP Profile for RTCP-based Feedback [9], and the Extended Secure RTP Profile for RTCP-based Feedback [10] MAY be used with DCCP (noting the potential conflict between DCCP partial checksums and the integrity protection provided by the secure RTP variants -- see Section 6).

5. RTP over DCCP: Signalling using SDP

The Session Description Protocol (SDP) [3] and the offer/answer model [11] are widely used to negotiate RTP sessions (for example, using the Session Initiation Protocol [22]). This section describes how SDP is used to signal RTP sessions running over DCCP.

5.1. Protocol Identification

SDP uses a media ("m=") line to convey details of the media format and transport protocol used. The ABNF syntax of a media line is as follows (from [3]):

```
media-field = %x6d "=" media SP port [ "/" integer ] SP proto
              1*(SP fmt) CRLF
```

The proto field denotes the transport protocol used for the media, while the port indicates the transport port to which the media is sent. Following [5] and [12] this memo defines the following five values of the proto field to indicate media transported using DCCP:

```
DCCP
DCCP/RTP/AVP
DCCP/RTP/SAVP
DCCP/RTP/AVPF
DCCP/RTP/SAVPF
```

The "DCCP" protocol identifier is similar to the "UDP" and "TCP"

protocol identifiers and describes the transport protocol, but not the upper-layer protocol. An SDP "m=" line that specifies the "DCCP" protocol MUST further qualify the application layer protocol using a "fmt" identifier (the "fmt" namespace is managed in the same manner as for the "UDP" protocol identifier). A single DCCP port is used, as denoted by the port field in the media line. The "DCCP" protocol identifier MUST NOT be used to signal RTP sessions running over DCCP; those sessions MUST use a protocol identifier of the form "DCCP/RTP/..." as described below.

The "DCCP/RTP/AVP" protocol identifier refers to RTP using the RTP Profile for Audio and Video Conferences with Minimal Control [4] running over DCCP.

The "DCCP/RTP/SAVP" protocol identifier refers to RTP using the Secure Real-time Transport Protocol [8] running over DCCP.

The "DCCP/RTP/AVPF" protocol identifier refers to RTP using the Extended RTP Profile for RTCP-based Feedback [9] running over DCCP.

The "DCCP/RTP/SAVPF" protocol identifier refers to RTP using the Extended Secure RTP Profile for RTCP-based Feedback [10] running over DCCP.

RTP payload formats used with the "DCCP/RTP/AVP", "DCCP/RTP/SAVP", "DCCP/RTP/AVPF" and "DCCP/RTP/SAVPF" protocol identifiers MUST use the payload type number as their "fmt" value. If the payload type number is dynamically assigned, an additional "rtpmap" attribute MUST be included to specify the format name and parameters as defined by the media type registration for the payload format.

Port 5004 is registered for use by RTP and SHOULD be the default port chosen by applications. If multiple RTP sessions are active from a host, even numbered ports in the dynamic range SHOULD be used for the other sessions. If RTCP is to be sent on a separate DCCP connection to RTP, the RTCP connection SHOULD use the next higher destination port number, unless an alternative DCCP port is signalled using the "a=rtcp:" attribute [13].

5.2. Service Codes

In addition to the port number, specified on the SDP "m=" line, a DCCP connection has an associated service code. A single new SDP attribute ("dccc-service-code") is defined to signal the DCCP service code according to the following ABNF [14]:

```
dccp-service-attr = %x61 "=dccp-service-code:" service-code
service-code      = hex-sc / decimal-sc / ascii-sc
hex-sc            = %x53 %x43 "=" %x78 *HEXDIG
decimal-sc       = %x53 %x43 "=" *DIGIT
ascii-sc         = %x53 %x43 ":" *sc-char
sc-char          = %d42-43 / %d45-47 / %d63-90 / %d95 / %d97-122
```

where DIGIT and HEXDIG are as defined in [14]. The service code is interpreted as defined in Section 8.1.2 of [2]. The following DCCP service codes are registered for use with RTP:

- o SC:RTPA an RTP session conveying audio data (and associated RTCP)
- o SC:RTPV an RTP session conveying video data (and associated RTCP)
- o SC:RTPT an RTP session conveying text media (and associated RTCP)
- o SC:RTPO an RTP session conveying other media (and associated RTCP)
- o SC:RTCP an RTCP connection, separate from the corresponding RTP

To ease the job of middleboxes, applications SHOULD use these service codes to identify RTP sessions running within DCCP.

The "a=dccp-service-code:" attribute is a media level attribute which is not subject to the charset attribute.

5.3. Connection Management

The "a=setup:" attribute indicates which of the end points should initiate the DCCP connection establishment (i.e. send the initial DCCP-Request packet). The "a=setup:" attribute MUST be used in a manner comparable with [12], except that DCCP connections are being initiated rather than TCP connections.

After the initial offer/answer exchange, the end points may decide to re-negotiate various parameters. The "a=connection:" attribute MUST be used in a manner compatible with [12] to decide whether a new DCCP connection needs to be established as a result of subsequent offer/answer exchanges, or if the existing connection should still be used.

5.4. Multiplexing Data and Control

A single DCCP connection can be used to transport multiplexed RTP and RTCP packets. Such multiplexing MUST be signalled using an "a=rtcp-mux" attribute according to [7]. If multiplexed RTP and RTCP is not to be used, then the "a=rtcp-mux" attribute MUST NOT be present in the SDP offer, and a separate DCCP connection MUST be opened to transport the RTCP data on a different DCCP port.

5.5. Example

An offerer at 192.0.2.47 signals its availability for an H.261 video session, using RTP/AVP over DCCP with service code "RTPV" (using the hexadecimal encoding of the service code in the SDP). RTP and RTCP packets are multiplexed onto a single DCCP connection:

```
v=0
o=alice 1129377363 1 IN IP4 192.0.2.47
s=-
c=IN IP4 192.0.2.47
t=0 0
m=video 5004 DCCP/RTP/AVP 99
a=rtcp-mux
a=rtpmap:99 h261/90000
a=dccp-service-code:SC=x52545056
a=setup:passive
a=connection:new
```

An answerer at 192.0.2.128 receives this offer and responds with the following answer:

```
v=0
o=bob 1129377364 1 IN IP4 192.0.2.128
s=-
c=IN IP4 192.0.2.128
t=0 0
m=video 9 DCCP/RTP/AVP 99
a=rtcp-mux
a=rtpmap:99 h261/90000
a=dccp-service-code:SC:RTPV
a=setup:active
a=connection:new
```

The end point at 192.0.2.128 then initiates a DCCP connection to port 5004 at 192.0.2.47. DCCP port 5004 is used for both the RTP and RTCP data, and port 5005 is unused. The textual encoding of the service code is used in the answer, and represents the same service code as in the offer.

6. Security Considerations

The security considerations in the RTP specification [1] and any applicable RTP profile (e.g. [4], [8], [9], or [10]) or payload format apply when transporting RTP over DCCP.

The security considerations in the DCCP specification [2] apply.

The SDP signalling described in Section 5 is subject to the security considerations of [3], [11], [12], [5], and [7].

The provision of effective congestion control for RTP through use of DCCP is expected to help reduce the potential for denial-of-service present when RTP flows ignore the advice in [1] to monitor packet loss and reduce their sending rate in the face of persistent congestion.

There is a potential conflict between the Secure RTP Profiles [8], [10] and the DCCP partial checksum option, since these profiles introduce, and recommend the use of, message authentication for RTP and RTCP packets. Message authentication codes of the type used by these profiles cannot be used with partial checksums, since any bit-error in the DCCP packet payload will cause the authentication check to fail. Accordingly, DCCP partial checksums SHOULD NOT be used in conjunction with SRTP authentication. The confidentiality features of the basic RTP specification cannot be used with DCCP partial checksums, since bit errors propagate. Also, despite the fact that bit errors do not propagate when using AES in counter mode, the Secure RTP profiles SHOULD NOT be used with DCCP partial checksums, since it requires authentication for security, and authentication is incompatible with partial checksums.

7. IANA Considerations

The following SDP "proto" field identifiers are registered: "DCCP", "DCCP/RTP/AVP", "DCCP/RTP/SAVP", "DCCP/RTP/AVPF" and "DCCP/RTP/SAVPF" (see Section 5.1 of this memo).

One new SDP attribute ("dccp-service-code") is registered (see Section 5.2 of this memo).

The following DCCP service codes are registered: SC:RTPA, SC:RTPV, SC:RTPT, SC:RTPO, and SC:RTCP (see Section 5.2 of this memo).

DCCP ports 5004 ("DCCP RTP") and 5005 ("DCCP RTCP") are registered for use as default RTP/RTCP ports (see Section 5.1 of this memo). The four services codes listed above are to be associated with these

ports.

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