



University  
of Glasgow

# The Session Layer and the DNS

Networked Systems 3  
Lecture 15

# Lecture Outline

- Higher layer protocols
- The session layer
  - Managing connections
  - Middleboxes and caches
  - Naming users and resources
- The domain name system (DNS)

# Higher Layer Protocols

- The OSI reference model defines three layers above the transport layer:
  - Session layer
  - Presentation layer
  - Application layer
- All typically implemented within an application or library; poorly-defined boundaries between layers

# Function of the Higher Layers

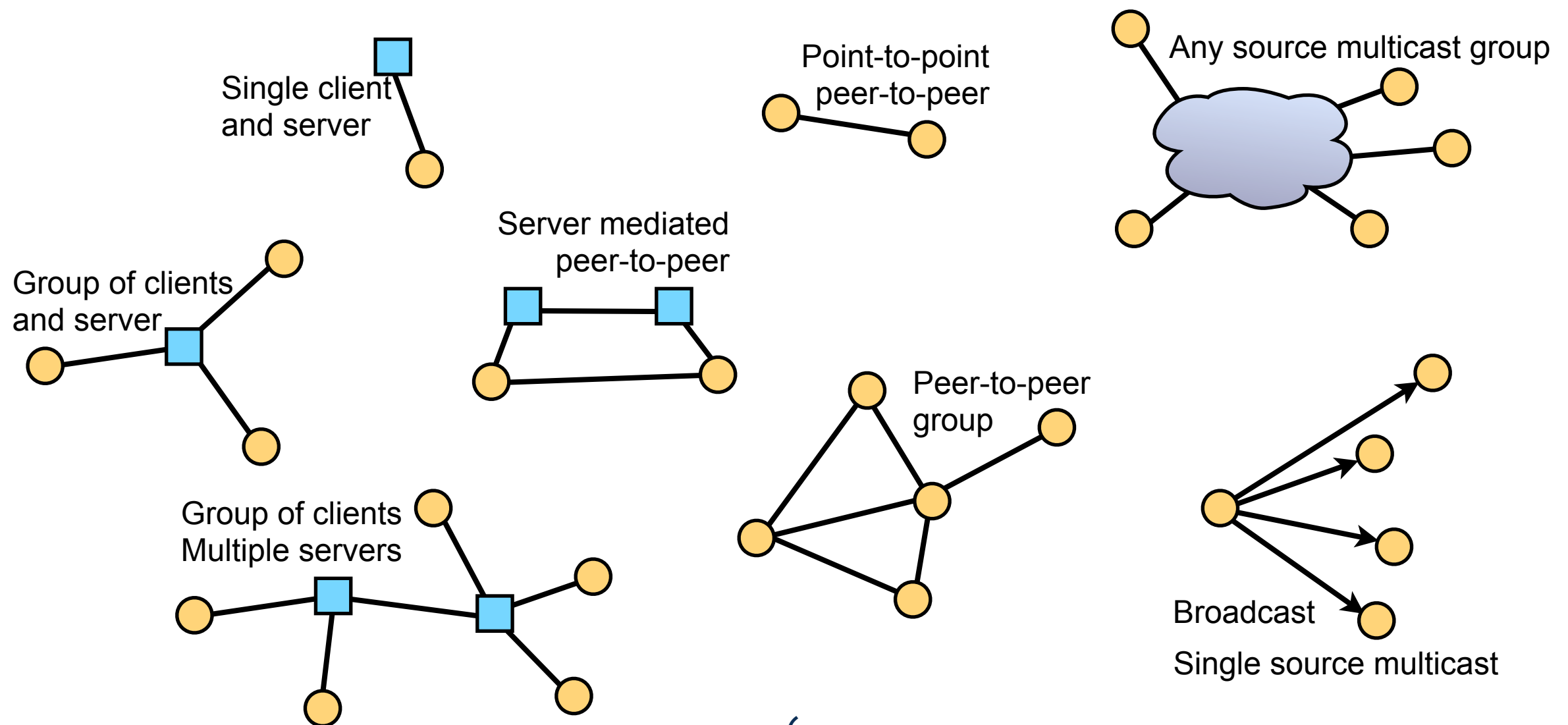
- To support the needs of the application:
  - Setup and manage transport layer connections
  - Name and locate application-level resources
  - Negotiate supported data formats, performing format conversion as needed
  - Present data in an appropriate manner
  - To implement application semantics

# The Session Layer

- Responsible for managing connections:
  - Find users/resources; create transport connections
  - Middleboxes and caches
- Responsible for naming resources:
  - Uniform resource identifiers
  - The Domain Name System (DNS)

# Managing Connections

- What connections does the application need?



# Managing Connections

- How to find participants?
  - Look-up name in a directory (e.g. DNS, web search engine)
  - Server mediated connection (e.g. instant messenger, VoIP call)
- How to setup connections?
  - Direct connection to named host (→ NAT issues)
  - Mediated service discovery, followed by peer-to-peer connection
    - E.g. VoIP using SIP and RTP with ICE
- How does session membership change?
  - Does the group size vary greatly? How rapidly do participants join and leave? Are all participants aware of other group members?

# User and Resource Mobility

- IP addresses encode location → mobility breaks transport layer connections
- Session layer must find new location, establish new connections
  - Old location might redirect – e.g., HTTP
  - Users might register new location
    - Updating a DNS name to point to the new IP address
    - Via an application-specific server – e.g., SIP proxy for VoIP calls



# Example: HTTP redirect

## HTTP request

```
GET /index.html HTTP/1.1  
Host: www.google.com
```

## HTTP response

```
HTTP/1.1 302 Moved Temporarily  
Location: http://www.google.co.uk/index.html  
Cache-Control: private  
Content-Type: text/html  
Server: gws  
Content-Length: 231  
Date: Sun, 17 Feb 2008 23:23:30 GMT  
  
<HTML><HEAD><meta http-equiv="content-type" content="text/html; charset=utf-8">  
<TITLE>302 Moved</TITLE></HEAD><BODY>  
<H1>302 Moved</H1>  
The document has moved  
<A HREF="http://www.google.co.uk/index.html">here</A>.  
</BODY></HTML>
```

302 response code indicates the content has moved, the “Location:” header specifies where it’s moved to.

# Multiple Connections

- A single session may span multiple transport connections
  - E.g., retrieving a web page containing images – one connection for the page, then one per image
  - E.g., a peer-to-peer file sharing application, building a distributed hash table
- Session layer responsible for co-ordinating the connections

# Middleboxes and Caches

- Some protocols rely on middleboxes or caches
  - Web cache – optimise performance, moving popular content closer to hosts
  - Email server – supports disconnected operation by holding mail until user connects
  - SIP proxy servers and instant messaging servers – locate users, respond for offline users
- The end-to-end argument applies, once again
  - Only add middleboxes when absolutely necessary

# How to Find the Middlebox?

- Manual configuration
- Look-up in central directory service
  - E.g., DNS MX records to find email servers
- Multicast service discovery
- “Transparent” redirection
  - E.g., Wi-Fi hotspots that grab web traffic, and redirect to a payment server

# Naming

- How to identify resources used or referenced by an application?
  - Files, email addresses, phone numbers, objects in a database, books, parcels being shipped, etc.
  - Use a *uniform resource identifier*
    - Uniform resource name (URN) – a unique resource name; no information on where to find, or how to access, the resource
    - Uniform resource locator (URL) – a unique resource name, plus location and access method
    - Directory service used for URN → URL mapping

# Uniform Resource Identifier

A general mechanism for naming arbitrary resources

**scheme**:**authority**/**path**?**query**#**fragment**

(authority, query and fragment optional)

**ftp**://**ftp.is.co.za/rfc/rfc1808.txt**

**http**://**news.bbc.co.uk/2/hi/europe/7249034.stm#map**

**ldap**://[**2001:db8::7**]/**c=GB?objectClass=one**

**mailto**:**John.Doe@example.com**

**news**:**comp.infosystems.www.servers.unix**

**tel**:**+1-816-555-1212**

**telnet**://**192.0.2.16:80/**

**urn**:**oasis:names:specification:docbook:dtd:xml:4.1.2**

Syntax is extremely flexible

Wide range of schemes defined

Some can be directly accessed, others require a look-up to map from the URI to a URL

# Domain Name System

- URIs often refer to a host on the network
  - Want to use a human-readable hostname in URIs, rather than an IP address
  - The *domain name system* (DNS) translates from the hostname to an IP address
    - www.dcs.gla.ac.uk → 130.209.240.1
    - DNS is an application layer protocol, running over the network
    - Not necessary for the correct operation of the transport or network layers, or lower

# History of the DNS

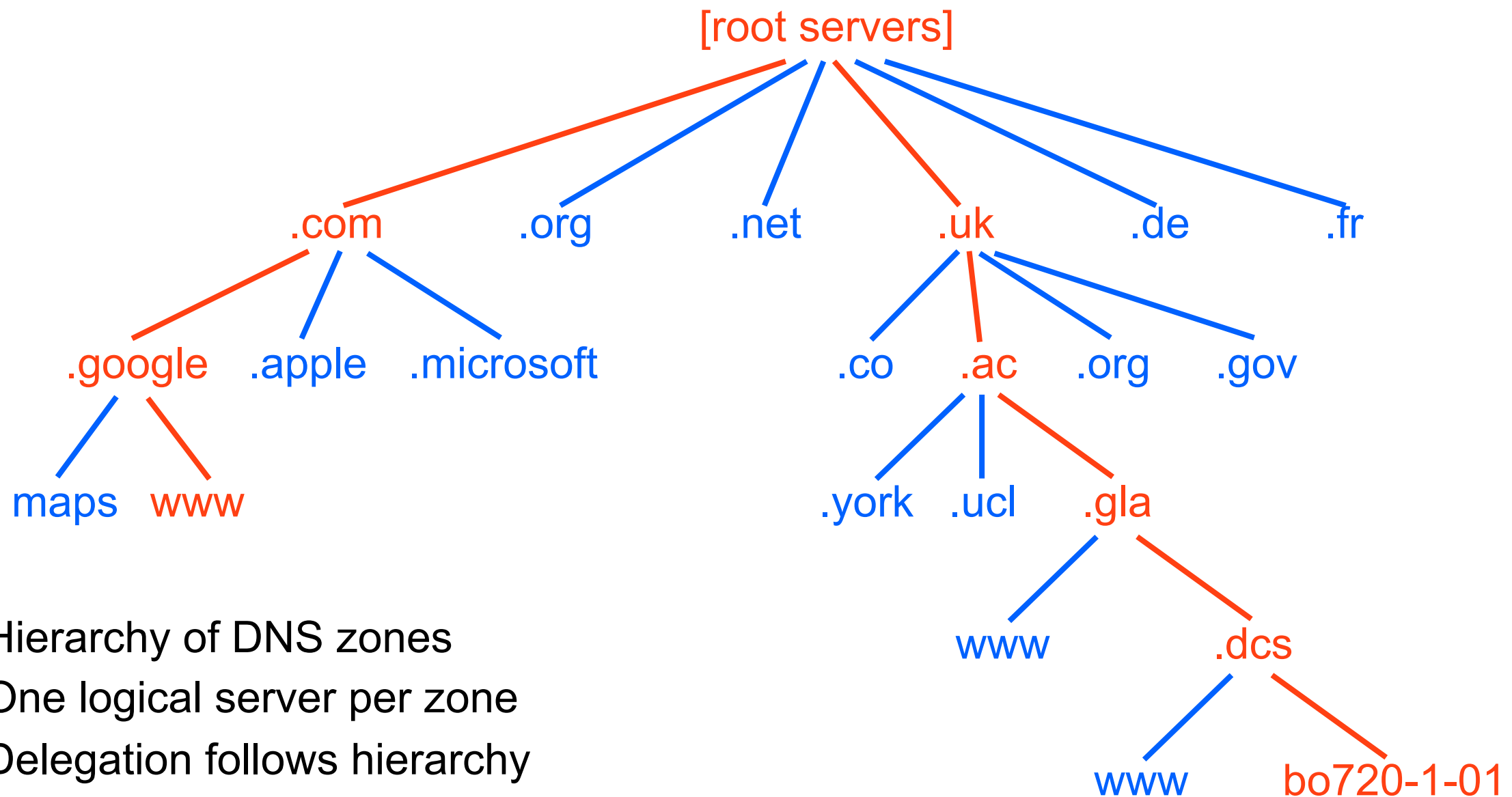
- Early Internet didn't use DNS
  - Flat file `hosts.txt` listing all host names and addresses
  - Maintained by central NIC; updated by email every few days; manually installed in hosts
- DNS proposed in 1983 as a distributed database of host names
  - Solve scaling problems with `hosts.txt`



Paul Mockapetris



# Operation of the DNS



- Hierarchy of DNS zones
- One logical server per zone
- Delegation follows hierarchy
- Hop-by-hop name look-up, follows hierarchy via root
- Results have TTL, cached at intermediate servers
- getaddrinfo()

# Contents of a DNS Zone

```
$TTL 3600          ; 1 hour
example.org.       IN      SOA      ns1.example.org. admin.example.org. (
                                2006051501      ; Serial
                                10800            ; Refresh
                                3600             ; Retry
                                604800          ; Expire
                                86400           ; Minimum TTL
                                )

; DNS Servers
                                IN      NS      ns1.example.org.
                                IN      NS      ns2.example.org.

; MX Records
                                IN      MX  10    mx.example.org.
                                IN      MX  20    mail.example.org.

; Machine Names
ns1      IN      A      192.168.1.2
ns2      IN      A      192.168.1.3
mx       IN      A      192.168.1.4
mail     IN      A      192.168.1.5
mail     IN      AAAA    2001:200:1000:0:25f:23ff:fe80:1234
server1  IN      A      192.168.1.10
server2  IN      A      192.168.1.11

; Aliases
www      IN      CNAME   server1
```

Source: adapted from The FreeBSD Handbook

# DNS Politics

- The DNS was administered by IANA
  - Jon Postel was IANA from its creation until his death in 1998
  - <http://www.ietf.org/rfc/rfc2468.txt> – “I remember IANA”
- DNS now managed into ICANN
  - The US government asserts ultimate control over ICANN, and hence the DNS
  - Significant attempts to move control of national domains to the UN, and hence to the countries concerned
  - Other attempts to set up alternate roots for the DNS, with different namespaces → significant technical problems
- Technical work ongoing to add non-ASCII domain names



Jon Postel

# Questions?